# Quantum Coherence Manipulation with Finite Resources

#### Kun Fang

PhD at UTS:QSI Sydney

Postdoc at DAMTP, Cambridge

Selected results from **1711.10512** (PRL) **1805.04045** (QUANTUM) + **1804.09500** (PRL)

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#### Talk Outline

©Resource Theory

Quantum Coherence Distillation How many cobits can be distilled?



OApplication: Channel Simulation How many cobits are required for simulation?



Summary and Discussions

## Resource Theory



#### Resource theory

#### Resource Theory = free states + free operations

Entanglement theory = separable states + LOCC

Steering = non-steerable + 1-way LOCC

Athermality = Gibbs states + thermal operations

Coherence theory = incoherent state + incoherent operations (next slide)

Like different currencies: CAD, EUR, GBP, ... trade between ...

[Chitambar-Gour-2019] RMP **1806.06107** 

Two perspectives of studying:

the asymptotic limits given infinite (i.i.d.) copies of resources?

tradeoff between param. given finite copies (one-shot) of resources?

#### Coherence theory

**Free states:** incoherent (diagonal) states  $\mathcal{I}:=\left\{ 
ho\geq 0: {
m Tr}\, 
ho=1, 
ho=\Delta(
ho)
ight\}$ 

Resource states: coherent (non-diagonal) states

diagonal map

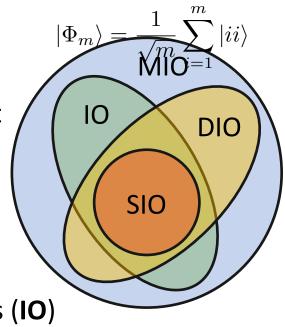
 $|\Phi_2\rangle$ 

Maximally coherent state:  $|\Psi_m\rangle=rac{1}{\sqrt{m}}\sum_{i=1}^m|i\rangle$  Coherent bit (cobit):  $|\Psi_2\rangle$ 

#### Free operations:

Maximally incoherent operations (MIO)

$$\rho \in \mathcal{I} \Longrightarrow \mathcal{E}(\rho) \in \mathcal{I}$$
$$[\Delta \circ \mathcal{E} \circ \Delta = \mathcal{E} \circ \Delta]$$



Dephasing-covariant incoherent operations (**DIO**)  $\mathcal{E} \circ \Delta = \Delta \circ \mathcal{E}$ 

ebit

Incoherent operations (IO)

$$\mathcal{E}(\cdot) = \sum_{i} E_{i} \cdot E_{i}^{\dagger},$$

$$E_{i} \cdot E_{i}^{\dagger} \in \text{MIO } \forall i$$

Strictly incoherent operations (**SIO**)

$$\mathcal{E}(\cdot) = \sum_{i} E_{i} \cdot E_{i}^{\dagger},$$
$$E_{i} \cdot E_{i}^{\dagger} \in \text{DIO } \forall i$$

#### Coherence theory

**Free states:** incoherent (diagonal) states  $\mathcal{I}:=\left\{ 
ho\geq0:\operatorname{Tr}
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ight\}$ 

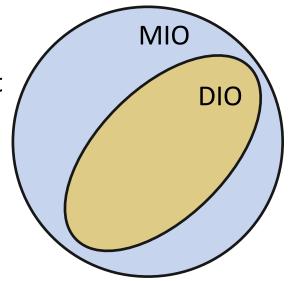
Resource states: coherent (non-diagonal) states

Maximally coherent state:  $|\Psi_m\rangle=rac{1}{\sqrt{m}}\sum_{i=1}^m|i
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#### Free operations:

Maximally incoherent operations (MIO)

$$\rho \in \mathcal{I} \Longrightarrow \mathcal{E}(\rho) \in \mathcal{I}$$
$$[\Delta \circ \mathcal{E} \circ \Delta = \mathcal{E} \circ \Delta]$$



Dephasing-covariant incoherent operations (**DIO**)

$$\mathcal{E} \circ \Delta = \Delta \circ \mathcal{E}$$

[Streltsov-Adesso-Plenio-2017] RMP 1609.02439

#### Why coherence?

Because we can indeed use coherence as a resource in some applications.

This talk: quantum channel simulation.

Other applications of coherence:

PHYSICAL REVIEW A 93, 012111 (2016)

Coherence as a resource in decision problems: The Deutsch-Jozsa algorithm and a variation

Mark Hillery

PRL 116, 240405 (2016)

PHYSICAL REVIEW LETTERS

week ending 17 JUNE 2016

#### **Entanglement and Coherence in Quantum State Merging**

A. Streltsov, 1,2,\* E. Chitambar, S. Rana, M. N. Bera, A. Winter, 4,5 and M. Lewenstein, 5

#### Quantum state redistribution with local coherence

Anurag Anshu,<sup>1</sup> Rahul Jain,<sup>2</sup> and Alexander Streltsov<sup>3,4</sup>

#### Why coherence?

#### ARTICLE

Received 20 Oct 2015 | Accepted 22 Apr 2016 | Published 20 May 2016

DOI: 10.1038/ncomms11712

**OPEN** 

# Numerical approach for unstructured quantum key distribution

Patrick J. Coles<sup>1</sup>, Eric M. Metodiev<sup>1</sup> & Norbert Lütkenhaus<sup>1</sup>

It is interesting to point out the connection to coherence<sup>44</sup>. For some set of orthogonal projectors  $\Pi = \{\Pi^j\}$  that decompose the identity,  $\sum_j \Pi^j = 1$ , the coherence (sometimes called relative entropy of coherence) of state  $\rho$  is defined as<sup>44</sup>:

$$\Phi(\rho, \Pi) = D\left(\rho \left\| \sum_{j} \Pi^{j} \rho \Pi^{j} \right). \tag{42}$$

Rewriting the primal problem in terms of coherence gives

$$\alpha = \min_{\rho_{AB} \in \mathcal{C}} \Phi(\rho_{AB}, Z_A). \tag{43}$$

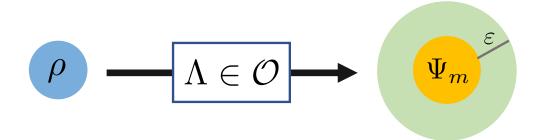
Hence, we make the connection that calculating the secret key rate is related to optimizing the coherence.

This observation is important since the coherence is a continuous function of  $\rho$  (Supplementary Note 6). This allows us to argue in Supplementary Note 6 that our optimization problem satisfies the strong duality criterion<sup>16</sup>, which means that the solution of the dual problem is precisely equal to that of primal problem.

### **Coherence Distillation**



#### One-shot coherence distillation



One-shot distillable coherence

$$C_{d,\mathcal{O}}^{(1),\varepsilon}(\rho) := \max_{\Lambda \in \mathcal{O}} \log_2 m$$
 s.t.  $F(\Lambda(\rho), \Psi_m) \ge 1 - \varepsilon$ .

The asymptotic distillable coherence

$$C_{d,\mathcal{O}}^{\infty}(\rho) = \lim_{\varepsilon \to 0} \lim_{n \to \infty} \frac{1}{n} C_{d,\mathcal{O}}^{(1),\varepsilon}(\rho^{\otimes n})$$

Question: can we compute  $C_{d,\mathcal{O}}^{(1),\varepsilon}(\rho)$  ?

#### Distillation under MIO and DIO

#### MIO

$$\begin{split} C_{d,\text{MIO}}^{(1),\varepsilon}(\rho) &= -\log_2 \min \eta \\ \text{s.t. } \text{Tr } G\rho &\geq 1-\varepsilon, \ 0 \leq G \leq \mathbb{1}, \ \Delta(G) = \eta \mathbb{1}. \end{split}$$

# $\min x + y$ s.t. $2x + 3y \le 1$ , $x - 5y \ge 4$ .

#### Semidefinite program (SDP)

- A generalization of linear program (LP) from scalar variables to linear operators
- © Efficiently computable
  - CVX for MATLAB
  - QI toolbox QETLAB by N. Johnston
- O Duality (Primal-Dual problems)
- Refer to J. Watrous' lecture note for details

#### Distillation under MIO and DIO

MIO

$$\begin{split} C_{d,\text{MIO}}^{(1),\varepsilon}(\rho) &= -\log_2\min\eta\\ \text{s.t. } \text{Tr } G\rho \geq 1-\varepsilon,\ 0 \leq G \leq \mathbb{1},\ \Delta(G) = \eta\mathbb{1}. \end{split}$$

#### Distillation under MIO and DIO

MIO

$$C_{d,\mathrm{MIO}}^{(1),arepsilon}(
ho) = -\log_2\min\eta$$
 s.t.  $\mathrm{Tr}\,G
ho \geq 1 - arepsilon,\ 0 \leq G \leq 1,\ \Delta(G) = \eta 1.$ 

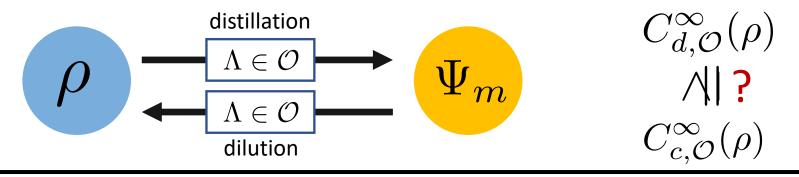
DIO

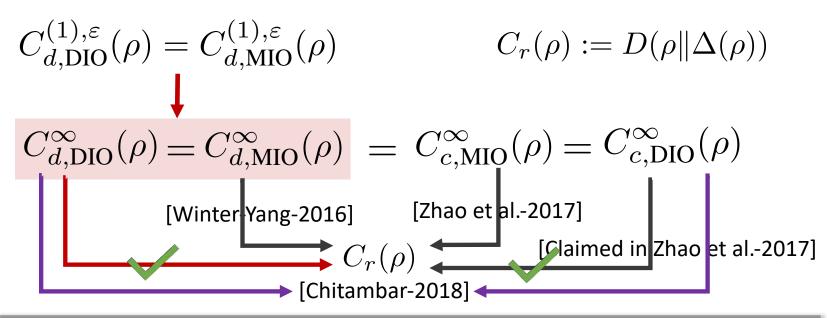
$$C_{d,\mathrm{DIO}}^{(1),arepsilon}(
ho) = -\log_2\min\eta$$
  
s.t.  $\mathrm{Tr}\,G
ho \geq 1 - arepsilon,\ 0 \leq G \leq \mathbb{1},\ \Delta(G) = \eta\mathbb{1}.$ 

Q: What is the difference between these two SDPs?

#### **NO** difference!

#### Asymptotic reversibility





**Reversibility** for entanglement theory [Brandão-Plenio-2010] and other resource theories [Brandão-Gour-2015] only known under **resource (asymptotically) non-generating maps**.

#### Hypothesis testing characterization

$$\mathcal{O} \in \{ MIO, DIO \}$$

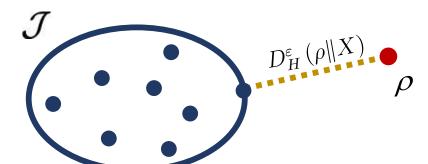
$$C_{d,\mathcal{O}}^{(1),\varepsilon}(\rho) = -\log_2 \min \frac{\eta}{\eta}$$
s.t. Tr  $G\rho \ge 1 - \varepsilon$ ,  $0 \le G \le 1$ ,  $\Delta(G) = \eta 1$ .

What does this SDP mean?

Hypothesis testing relative entropy:

$$\begin{split} D_H^\varepsilon(\rho\|\sigma) := -\log_2 \min \mathrm{Tr}\, G\rho \\ \mathrm{s.t.}\, \mathrm{Tr}\, G\sigma \geq 1 - \varepsilon, \ 0 \leq G \leq \mathbb{1}. \end{split}$$

$$C_{d,\mathrm{MIO}}^{(1),\varepsilon}(\rho) = \min_{X \in \mathcal{J}} D_H^{\varepsilon}(\rho \| X) \quad \mathrm{with} \quad \mathcal{J} = \{X : \mathrm{Tr}\, X = 1, \Delta(X) = X\}$$



Not necessarily positive semidefinite!

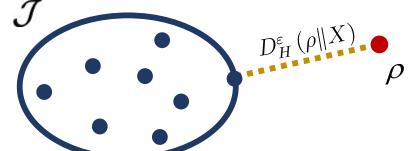
#### Comparison

Compare with other similar results e.g., [Zhao-Liu-Yuan-Chitambar-Winter-2019]

$$\min_{\sigma \in \mathcal{I}} D_{\min}^{\frac{\varepsilon}{2} - \eta}(\rho \| \sigma) + 2\log \eta \le C_{d, \text{IO}}^{(1), \varepsilon}(\rho) \le \min_{\sigma \in \mathcal{I}} D_{\min}^{\sqrt{\varepsilon(2 - \varepsilon)}}(\rho \| \sigma)$$

bounds + correction terms

$$C_{d,\mathrm{MIO}}^{(1),\varepsilon}(\rho) = \min_{X \in \mathcal{J}} D_H^{\varepsilon}(\rho \| X) \quad \mathrm{with} \quad \mathcal{J} = \{X : \mathrm{Tr}\, X = 1, \Delta(X) = X\}$$



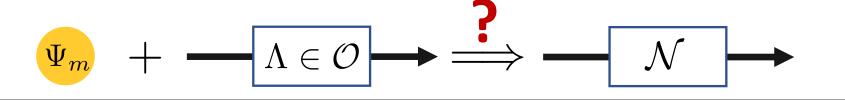
Not necessarily positive semidefinite!

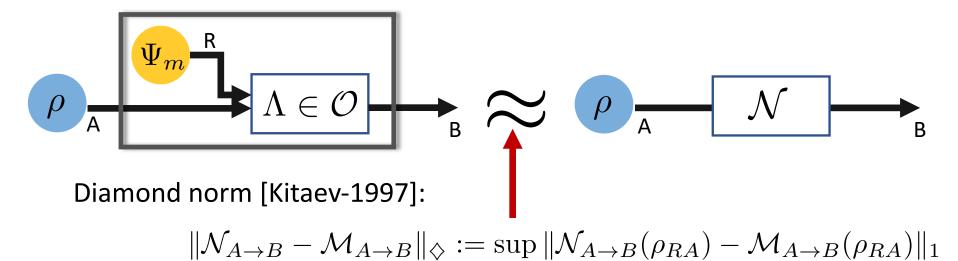
weird + clean

# Channel Simulation via Quantum Coherence



#### Channel simulation setting





One-shot simulation cost

$$S_{c,\mathcal{O}}^{(1),\varepsilon} := \min_{\Lambda \in \mathcal{O}} \log_2 m \quad \text{s.t. } \frac{1}{2} \|\Lambda_{AR \to B}(\Psi_m \otimes \cdot) - \mathcal{N}_{A \to B}(\cdot)\|_{\diamondsuit} \leq \varepsilon$$

 $\rho_{RA}$ 

Question: can we compute  $S_{c,\mathcal{O}}^{(1),\varepsilon}$  ?

#### SDP for simulation cost

#### SDP for diamond norm [Watrous-2009]:

$$\frac{1}{2} \|\mathcal{N}_1 - \mathcal{N}_2\|_{\diamondsuit} = \min \gamma$$
s.t.  $\operatorname{Tr}_B Y_{AB} \le \gamma \cdot \mathbb{1}_A$ ,
$$Y_{AB} \ge J_{\mathcal{N}_1} - J_{\mathcal{N}_2}, \ Y_{AB} \ge 0.$$

$$S_{c,\text{MIO}}^{(1),\varepsilon}(\mathcal{N}) = \log_2 \min \operatorname{Tr} J_{\widetilde{\mathcal{M}}} / d_A$$

$$\text{s.t. } 0 \leq J_{\widetilde{\mathcal{N}}} \leq J_{\widetilde{\mathcal{M}}},$$

$$\operatorname{Tr}_B J_{\widetilde{\mathcal{N}}} = \mathbb{1}_A,$$

$$\operatorname{Tr}_B J_{\widetilde{\mathcal{M}}} = \operatorname{Tr} J_{\widetilde{\mathcal{M}}} / d_A \cdot \mathbb{1}_A,$$

$$\operatorname{Tr}_A J_{\widetilde{\mathcal{M}}} (|i\rangle\langle i| \otimes \mathbb{1}_B) \in \Delta, \ \forall i$$

$$\operatorname{Tr}_B Y_{AB} \leq \varepsilon \cdot \mathbb{1}_A$$

$$Y_{AB} \geq J_{\widetilde{\mathcal{N}}} - J_{\mathcal{N}}, Y_{AB} \geq 0.$$

Does not look nice but indeed an SDP!

The proof technique does not work for DIO.  $\Lambda_{RA\to B}$ 

#### Distance characterization

The one-shot coherence simulation cost under MIO is given by

$$S_{c,\text{MIO}}^{(1),\varepsilon}(\mathcal{N}) = \min_{\mathcal{M} \in \text{MIO}} D_{\max}^{\varepsilon}(\mathcal{N}||\mathcal{M}),$$

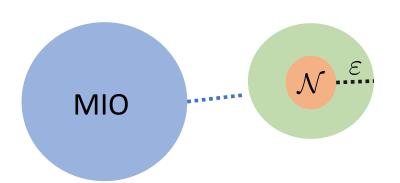
with the channel divergence

$$D_{\max}(\mathcal{N}||\mathcal{M}) := D_{\max}(J_{\mathcal{N}}||J_{\mathcal{M}}) \quad \text{and} \quad D_{\max}^{\varepsilon}(\mathcal{N}||\mathcal{M}) := \inf_{\substack{\frac{1}{2} \, ||\widehat{\mathcal{N}} - \mathcal{N}||_{\diamondsuit} \leq \varepsilon \\ \widehat{\mathcal{N}} \in \operatorname{CPTP}}} D_{\max}(\widehat{\mathcal{N}}||\mathcal{M}).$$

Compatible with other channel divergence

[Cooney-Mosonyi-Wilde-2016; Leditzky-Kaur-Datta-Wilde-2018]

$$\mathbf{D}(\mathcal{N}||\mathcal{M}) := \max_{|\varphi\rangle_{RA}} \mathbf{D}(\mathcal{N}_{A\to B}(\varphi_{RA})||\mathcal{M}_{A\to B}(\varphi_{RA}))$$



#### **Channel resource theory**

[KF-Wang-Tomomichel-Berta-2018] 1807.05354

[Faist-Berta-Brandão-2018] 1807.05610

[Gour-Wilde-2018] 1808.06980

[Gour-2018] 1808.02607

[Li-Bu-Liu-2018] 1812.02572

....

## Summary & Discussions

#### Summary

Quantum coherence distillation

How many cobits can be distilled?

SDP and entropy characterizations for one-shot distillable coherence under MIO and DIO which have the same power;

Application: Channel simulation

How many cobits are required for simulation?

SDP and entropy characterizations for the coherence cost of channel simulation under MIO.

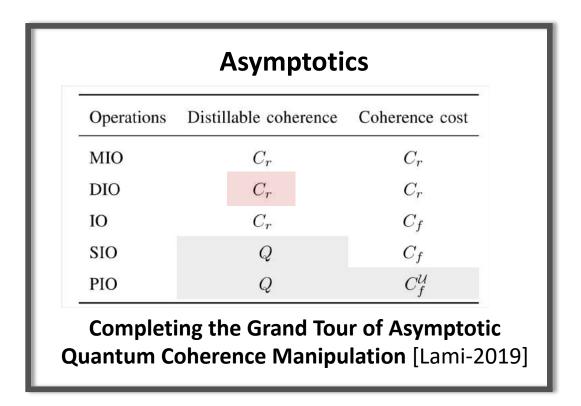
#### **One-shot** MIO DIO IO SIO $\widetilde{C}_H^{\varepsilon}$ [25] $\mathbf{C}_{\min}^{arepsilon'}$ [\*] distillation $\tilde{C}_H^{\varepsilon}$ [25] Thm. 10 [\*] One-shot formation $C_0^{\varepsilon}$ [24] $C_{\max}^{\varepsilon}$ [24] $C_{\Delta,\max}^{\varepsilon}$ [24] $C_0^{\varepsilon}$ [24]

One-Shot Coherence Distillation: Towards Completing the Picture [Zhao-Liu-Yuan-Chitambar-Winter-2018]

What about n-shot scenario?  $\rho^{\otimes n}$ 

Size of  $ho^{\otimes n}$  increases exponentially fast w.r.t. n







For example,

$$C_{d,\text{MIO}}^{(1),\varepsilon}(\rho^{\otimes n}) \stackrel{?}{=} nD(\rho \| \Delta(\rho)) + \sqrt{nV(\rho \| \Delta(\rho))} \Phi^{-1}(\varepsilon) + O(\log n)$$

# one-shot large blocklength Asympt.

One-shot simulation cost:

$$S_{c,\text{MIO}}^{(1),\varepsilon}(\mathcal{N}) = \min_{\mathcal{M} \in \text{MIO}} D_{\max}^{\varepsilon}(\mathcal{N}||\mathcal{M}),$$

Asymptotic simulation cost:

$$\lim_{\varepsilon \to 0} \lim_{n \to \infty} \frac{1}{n} S_{c,\text{MIO}}^{(1),\varepsilon}(\mathcal{N}^{\otimes n}) = \lim_{\varepsilon \to 0} \lim_{n \to \infty} \frac{1}{n} \min_{\mathcal{M}^n \in \text{MIO}} D_{\max}^{\varepsilon}(\mathcal{N}^{\otimes n} || \mathcal{M}^n) \longrightarrow ?$$

Partial progress: for classical-quantum channels

$$\longrightarrow \max_{\rho} C_r(\mathcal{N}(\rho)), \quad \text{where} \quad C_r(\rho) := D(\rho \| \Delta(\rho))$$

#### More about me

#### Webpage: kunfang.info

#### Also work on:

- o entanglement theory
- © channel capacity

#### Recently working on:

- polynomial optimization
- sum-of-square hierarchies
- uncertainty relations

entanglement-assisted
quantum communication
quantum coherence success probability
semidefinite programming
no-signaling quantum channel fidelity

distillation
single qubit LOCC one-shot broadcast resource states
coherent states classical communication
entanglement measure pure states entanglement
efficient computation
amplitude damping quantum correlations
quantum operations

(Generated by scimeter.org)

Happy to discuss any QUANTUM problems with or without optimizations!

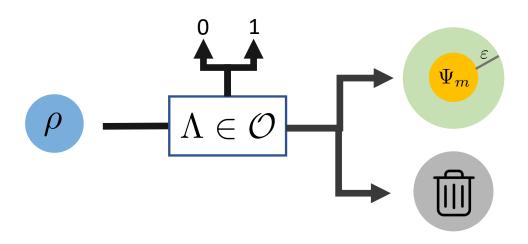
[IQC office: QNC 0213]

# Thanks for your attention!

This talk: 1711.10512 (PRL) 1805.04045 (QUANTUM) + 1804.09500 (PRL)

## Something more...

#### Probabilistic distillation



Maximum success probability of coherence distillation:

$$P_{\mathcal{O}}(\rho \to \Psi_m, \varepsilon) := \max p$$
s.t.  $\Lambda_{A \to FB}(\rho) = p|0\rangle\langle 0|_F \otimes \rho' + (1-p)|1\rangle\langle 1|_F \otimes \omega,$ 

$$F(\rho', \Psi_m) \ge 1 - \varepsilon, \Lambda \in \mathcal{O}.$$

Question: how to compute  $P_{\mathcal{O}}(\rho \rightarrow \Psi_m, \varepsilon)$  ?

#### SDP characterization

The maximal success probability of distillation under MIO/DIO are given by SDPs:

$$P_{\text{MIO}}(\rho \to \Psi_m, \varepsilon) = \max \operatorname{Tr} G \rho$$
  
s.t.  $\Delta(G) = m \Delta(C),$   
 $0 \le C \le G \le \mathbb{1},$   
 $\operatorname{Tr} C \rho \ge (1 - \varepsilon) \operatorname{Tr} G \rho.$ 

$$P_{\mathrm{DIO}}(\rho \to \Psi_m, \varepsilon) = \max \ \mathrm{Tr} \, G \rho$$
 s.t.  $\Delta(G) = m \Delta(C),$   $0 \le C \le G \le 1,$   $\mathrm{Tr} \, C \rho \ge (1 - \varepsilon) \, \mathrm{Tr} \, G \rho,$   $G = \Delta(G).$ 

#### No-go theorem

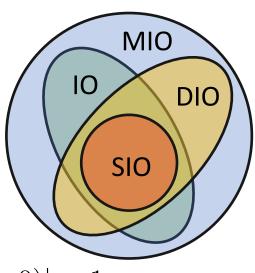
For any full-rank quantum state ho ,  $\$ it holds that  $P_{ ext{MIO}}(
ho o \Psi_m,0) = 0$  .

- MIO sets a fundamental limits for all other incoherent operations;
- Any generic density matrix has full rank;
- Open Depolarizing noise:

$$\mathcal{D}_{\alpha}(\rho) = (1 - \alpha) \cdot \rho + \alpha \cdot 1/m$$

 $\bigcirc$  Non-continuity:  $F(\mathcal{D}_{\varepsilon}(\Psi_m), \Psi_m) \to 1$ 

$$\mathsf{but} \left| P_{\mathsf{MIO}}(\mathcal{D}_{\varepsilon}(\Psi_{m}) \to \Psi_{m}, 0) - P_{\mathsf{MIO}}(\Psi_{m} \to \Psi_{m}, 0) \right| = 1$$



#### Pure state + MIO

For any pure states 
$$|\varphi\rangle=\sum_{i=1}^n \varphi_i|i\rangle,\; \varphi_i\neq 0,\; n\geq 2,\; \text{it holds}$$

$$P_{\text{MIO}}(\varphi \to \Psi_m, 0) \ge \frac{n^2}{m(\sum_{i=1}^n |\varphi_i|^{-2})} > 0$$

Use one cobit resource to win a million cobits lottery!



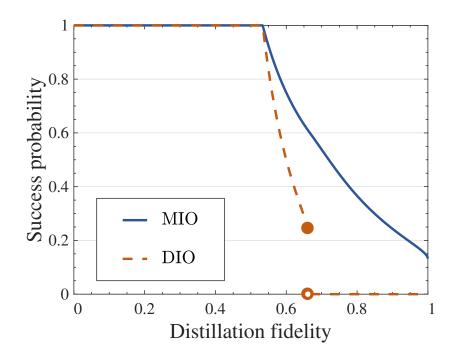
$$P_{\text{MIO}}(\Psi_2 \to \Psi_2^{\otimes 10^6}, 0) \ge \frac{1}{2^{10^6}}.$$

Not happening for DIO!

#### "Sudden death" for DIO

For any pure states 
$$|\varphi\rangle=\sum_{i=1}^n \varphi_i|i\rangle,\; \varphi_i\neq 0,\; n\geq 2,\; \text{it holds}$$

$$P_{\mathrm{DIO}}(\varphi \to \Psi_m, \varepsilon) \begin{cases} > 0 & \text{if } n \geq m \text{ or if } n < m \text{ and } \varepsilon \geq 1 - \frac{n}{m}, \\ = 0 & \text{if } n < m \text{ and } \varepsilon < 1 - \frac{n}{m}. \end{cases}$$



$$(|0\rangle + 3|1\rangle)/\sqrt{10} \rightarrow \Psi_3$$

``Analogous" to (pretty) strong converse theorem in channel coding theory